# Uncountably many well-quasi-ordered permutation classes 

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Joint work with Vince Vatter (U. Florida)

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Given a class $\mathcal{C}$, what properties guarantee a 'tame' enumeration?

## Finitely many simple permutations

## Theorem (Albert and Atkinson 2005)

If a class $\mathcal{C}$ contains only finitely many simple permutations, then it has an algebraic generating function and is finitely based.


## Geometrically griddable classes

Theorem (Albert, Atkinson, Bouvel, Ruškuc and Vatter 2013)
If a class $\mathfrak{C}$ is geometrically griddable, then it has a rational generating function and is finitely based.

regular language over finite alphabet

## Non-example: two stacks in series



Pierrot \& Rossin (2017) Membership is polynomial time

Elvey Price \& Guttman (2017) Exact enumeration to length 20 Generating function $\sim A(1-\mu \cdot z)^{\gamma}$

Murphy (2003) Not finitely based

## Finitely based classes

All classes that have finitely many simples, or that are geometrically griddable are finitely based. Two-stacks are not.

Conjecture (Noonan, Zeilberger, 1996)
Every finitely based class has a D-finite generating function.

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Theorem (Garrabrant, Pak, 2015)
Zeilberger is right: Noonan-Zeilberger is false.

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## Conjecture

Every finitely based class with growth rate $<4$ has a rational generating function.

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## Subclasses of $\operatorname{Av}(231), \operatorname{Av}(321)$

|  | $\mathcal{C} \subsetneq \operatorname{Av}(231)$ | $\mathcal{D} \subsetneq \operatorname{Av}(321)$ |
| :--- | :---: | :---: |
| Growth rate | Countably many <br> possibilities | Includes $[2.36,2.48]$ <br> (Bevan, 2018) |

Generating function

Basis

Rational (Albert, Atkinson, 2005)

Could be anything

Finite
Finite or infinite

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Infinite antichains?

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Every WQO permutation class has an algebraic generating function.

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# $\begin{array}{llllllllllll}0 & 1 & 1 & 0 & 1 & 0 & 0 & 0 & 1 & 1 & 0 & \cdots\end{array}$ 

Binary word $w \longrightarrow$ permutation $\pi_{w}$

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\begin{array}{lllllllllllll}
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\mathcal{P}=\left\{\text { permutations } \pi \text { contained in } \pi_{w_{i}} \text { for some } i\right\}
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Prouhet-Thue-Morse is uniformly recurrent $\Longrightarrow \mathcal{P}$ is WQO.

Prouhet-Thue-Morse ' $(1,1,1, \ldots)^{\prime}$

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\begin{aligned}
f_{\mathcal{P}}(z)= & 1+z+2 z^{2}+6 z^{3}+22 z^{4} \\
& +80 z^{5}+276 z^{6}+948 z^{7} \\
& +3276 z^{8}+\cdots
\end{aligned}
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Sequence ' $(2,1,1,1, \ldots$ )'
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f_{\mathfrak{Q}}(z)= & 1+z+2 z^{2}+6 z^{3}+22 z^{4} \\
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- The binary sequences have different complexity functions
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## Corollary

There exist WQO permutation classes that do not have algebraic (or even $D$-finite) generating functions.

## Concluding remarks

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In search of 'tame' enumeration:

- Is labelled WQO enough to guarantee algebraic g.f.s?
- (Note: $\mathrm{LWQO} \Longrightarrow \mathrm{WQO}+$ finitely based.)

Merci!

