# Fishburn Trees 

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## The Fishburn numbers

$$
\sum_{n \geq 0} \prod_{k=1}^{n}\left(1-(1-x)^{k}\right)=1+x+2 x^{2}+5 x^{3}+15 x^{4}+53 x^{5}+217 x^{6}+\ldots
$$

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## Milestone Papers

(1) (2+2)-free posets, ascent sequences and pattern avoiding permutations,
M. Bousquet-Melou, A. Claesson, M. Dukes, S. Kitaev, 2010.
-Ascent sequences (and modified ascent sequences);
-Fishburn permutations;
-Unlabeled (2+2)-free posets;
-Stoimenow matchings.

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## Milestone Papers

(ii) Ascent sequences and upper triangular matrices containing non-negative integers,
M. Dukes, R. Parviainen, 2010.
-Fishburn matrices and ascent sequences.

## The Fishburn numbers

$\sum_{n \geq 0} \prod_{k=1}^{n}\left(1-(1-x)^{k}\right)=1+x+2 x^{2}+5 x^{3}+15 x^{4}+53 x^{5}+217 x^{6}+\ldots$

## Milestone Papers

(III) Composition matrices, (2+2)-free posets and their specializations, M. Dukes, V. Jelínek, M. Kubitzke, 2011.
-Fishburn matrices and $(\mathbf{2}+\mathbf{2})$-free posets.

## The Fishburn numbers

$\sum_{n \geq 0} \prod_{k=1}^{n}\left(1-(1-x)^{k}\right)=1+x+2 x^{2}+5 x^{3}+15 x^{4}+53 x^{5}+217 x^{6}+\ldots$

## ONE MORE REFERENCE

(IV) Transport of patterns by Burge transpose, C., Claesson, PP2022.
-Modified ascent sequences and Fishburn permutations;
-Transport of patterns based on the Burge transpose.



## Modified Ascent Sequences

## Combinatorial definition (C., Clabsson, 2020)

A sequence $x=x_{1} \cdots x_{n}$ is a MODIFIED (ASCENT) SEQUENCE if:
(1) $x_{1}=1$;
(2) $\left\{x_{1}, \ldots, x_{n}\right\}$ is an interval;
(3) $x_{i}<x_{i+1}$ if and only if $x_{i+1}$ is the leftmost copy of $x_{i+1}$ in $x$.

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$$
\left\{x_{1}\right\} \cup\{\text { Ascent tops }\}=\{\text { Leftmost copies }\}
$$

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Cayley permutations avoiding the Cayley-mesh patterns

and


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## ExAMPLE

$$
\mathrm{x}=\underline{1} 1 \underline{4} 111 \underline{2} 2 \underline{5} 22 \underline{3}
$$

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## NON-EXAMPLE

$$
\mathrm{x}=1212
$$

## Modified Ascent Sequences

## Combinatorial definition (C., Clabsson, 2020)

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## ANOTHER NON-EXAMPLE: WHY?

$$
x=12143
$$

## From modified sequences to Fishburn trees

- Let $x_{m}$ be the leftmost copy of $\max (x)$ in $x$;
- The MAX-DECOMPOSITION of $x=x_{1} \ldots x_{n}$ is

$$
x=\operatorname{pref}(x) x_{m} \operatorname{suff}(x)
$$

## From modified sequences to Fishburn trees

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$$

The rooted, labeled, binary tree $\mathrm{T}(x)$ is defined by $\mathrm{T}(\emptyset)=\emptyset$ and:


$$
x=\quad \begin{array}{llllllllllllll} 
& 1 & 1 & 4 & 1 & 1 & 1 & 2 & 2 & 5 & 2 & 2 & 3
\end{array}
$$

$\mathrm{T}(x)=$
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$$
x=\begin{array}{llllllllllllll} 
& & 1 & 1 & 4 & 1 & 1 & 1 & 2 & 2 & 5 & 2 & 2 & 3
\end{array}
$$

$\mathrm{T}(x)=\quad 1$


$$
x=\begin{array}{llllllllllllll} 
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x=\begin{array}{llllllllllllll}
x= & 1 & 1 & 4 & 1 & 1 & 1 & 2 & 2 & 5 & 2 & 2 & 3
\end{array}
$$



Can you guess how to determine $x$ from the tree $\mathrm{T}(x)$ ?

$$
x=
$$

$\begin{array}{llllllllllll}1 & 1 & 4 & 1 & 1 & 1 & 2 & 2 & 5 & 2 & 2 & 3\end{array}$


Can you guess how to determine $x$ from the tree $\mathrm{T}(x)$ ? We use the in-order traversal!

Let $T=(L, r, R)$ be the rooted, labeled, binary tree with:

- Left subtree $L$;
- Root $r$ with label $\mathfrak{l}(r)$;
- Right subtree $R$


## IN-ORDER SEQUENCE

The in-order sequence of $T$ is defined recursively by:

$$
x(\emptyset)=\emptyset \quad \text { and } \quad x(T)=x(L) \mathfrak{l}(r) x(R)
$$

"In-order traverse $T$ and write down the label of each visited node".


$$
\mathrm{T}(x)=
$$

$$
x=\quad \begin{array}{llllllllllllll} 
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$$

Properties of $\mathrm{T}(x)$
(1) ROOTED.

## Properties of $T(x)$

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(2) BINARY: Each node has either:

- 0 children;
- 1 child, which is either the left or the right child;
- 2 children, namely a left child and a right child.


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© Rooted.
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(6) STRICTLY DECREASING TO THE LEFT.


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A Fishburn tree is a tree that satisfies all the above properties.

## MAXIMAL-RIGHT-PATH DECOMPOSITION

The diagonal of $T$ is the path from the root to the leftmost node.

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The diagonal of $T$ is the path from the root to the leftmost node.

- Decompose $T$ in maximal right paths.
- In each path, the topmost node $v$ is either:
(I) on the diagonal;
(ii) or, the left child of a node that is not on the diagonal.
- Number each maximal right path by the label of:
(I) its top node, if the top node is on the diagonal;
(iI) the father of its top node, otherwise.








## MAXIMAL-RIGHT-PATH DECOMPOSITION

## THEOREM

Let $k$ be the maximal value of a label in $T$.

- There are $k$ maximal-right-paths.
- Each one corresponds to a unique number in $\{1,2, \ldots, k\}$.


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EACH NODE $v$ OF $T$ GETS TWO LABELS:
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## EACH NODE $v$ OF $T$ GETS TWO LABELS:

The node label $\mathfrak{l}(v) ; \quad$ The path label $\mathfrak{b}(v)$.

## Theorem

$$
\mathfrak{l}(v) \leq \mathfrak{b}(u)
$$

Proof. A Fishburn tree is decreasing.

## Fishburn matrices

- Lower triangular;
- Non-negative integer entries;
- Every row and column contains at least one non-zero entry.
- The size of a Fishburn matrix is the sum of its entries.


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- The size of a Fishburn matrix is the sum of its entries.


## Fishburn matrices of size 3

$$
[3],\left[\begin{array}{ll}
2 & 1 \\
\cdot & 1
\end{array}\right],\left[\begin{array}{ll}
1 & 2 \\
\cdot & 2
\end{array}\right],\left[\begin{array}{ll}
1 & \\
1 & 1
\end{array}\right],\left[\begin{array}{lll}
1 & & \\
\cdot & 1 & \\
\cdot & \cdot & 1
\end{array}\right]
$$

The Fishburn matrix $A=\left(a_{i j}\right)$ associated with $T$ is:

$$
\begin{aligned}
a_{i j} & =\mid\{\text { nodes with label } j \text { contained in the } i \text { th path of } T\} \mid \\
& =|\{v: \mathfrak{l}(v)=i, \mathfrak{b}(v)=j\}| .
\end{aligned}
$$



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& =|\{v: \mathfrak{l}(v)=i, \mathfrak{b}(v)=j\}| .
\end{aligned}
$$



## From Fishburn matrices to Fishburn trees

$$
\left[\begin{array}{ccccc}
2 & & & & \\
3 & \cdot & & & \\
\cdot & 2 & \cdot & & \\
\cdot & 2 & \cdot & 1 & \\
\cdot & \cdot & 1 & \cdot & 1
\end{array}\right]
$$

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$$
\left[\begin{array}{ccccc}
2 & & & & \\
3 & \cdot & & & \\
\cdot & 2 & \cdot & & \\
\cdot & 2 & \cdot & 1 & \\
\cdot & \cdot & 1 & \cdot & 1
\end{array}\right] \quad \begin{aligned}
& 1 \\
& 1
\end{aligned}
$$

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$$
\left[\begin{array}{ccccc}
2 & & & & \\
3 & \cdot & & & \\
\cdot & 2 & \cdot & & \\
\cdot & 2 & \cdot & 1 & \\
\cdot & \cdot & 1 & \cdot & 1
\end{array}\right] \quad{ }_{1}^{1}
$$

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$$
\left[\begin{array}{ccccc}
2 & & & & \\
3 & \cdot & & & \\
\cdot & 2 & \cdot & & \\
\cdot & 2 & \cdot & 1 & \\
\cdot & \cdot & 1 & \cdot & 1
\end{array}\right] \quad \begin{aligned}
& 2 \\
&
\end{aligned} \quad \begin{array}{|l} 
\\
\end{array}
$$

## From Fishburn matrices to Fishburn trees

$$
\left[\begin{array}{lllll}
2 & & & & \\
3 & \cdot & & & \\
\cdot & 2 & \cdot & & \\
\cdot & 2 & \cdot & 1 & \\
\cdot & \cdot & 1 & \cdot & 1
\end{array}\right] \quad{ }_{2}^{4}
$$

## From Fishburn matrices to Fishburn trees

$$
\left[\begin{array}{lllll}
2 & & & & \\
3 & \cdot & & & \\
\cdot & 2 & \cdot & \\
\cdot & 2 & \cdot & 1 \\
\cdot & \cdot & 1 & \cdot & 1
\end{array}\right] \quad{ }^{5}
$$

## From Fishburn matrices to Fishburn trees

$$
\left[\begin{array}{llll}
\left.\begin{array}{lll}
2 & & \\
3 & \cdot & \\
\cdot & 2 & \cdot \\
\cdot & 2 & \cdot \\
\cdot & \cdot & 1
\end{array}\right] & { }^{1} & \\
& \mathbf{1}^{1} & (4)_{2}^{4} & \mathbf{5}^{3}
\end{array}\right.
$$

## From Fishburn matrices to Fishburn trees



## From Fishburn matrices to Fishburn trees

$$
\left[\begin{array}{cccc}
2 & & & \\
3 & \cdot & & \\
\cdot & 2 & \cdot & \\
\cdot & 2 & \cdot & 1 \\
\cdot & \cdot & 1 & \cdot
\end{array}\right]
$$



## From Fishburn matrices to Fishburn trees

$$
\left[\begin{array}{cccc}
\left.\begin{array}{cccc}
2 & & & \\
3 & \cdot & & \\
\cdot & 2 & \cdot & \\
\cdot & 2 & \cdot & 1 \\
\cdot & \cdot & 1 & \cdot
\end{array}\right]
\end{array}\right]
$$



## From Fishburn matrices to Fishburn trees

$$
\left[\begin{array}{cccc}
2 & & & \\
3 & \cdot & & \\
\cdot & 2 & \cdot & \\
\cdot & 2 & \cdot & 1 \\
\cdot & \cdot & 1 & \cdot
\end{array}\right]
$$



## "Embedding" of Trees on binary matrices

1

11

11

1 . . 1
. . 1 . 1

## "EMBEDDING" OF TREES ON BINARY MATRICES

(1)
(1)
$(1)$
(1) • (1)
(1) (1)

## "EMBEDDING" OF TREES ON BINARY MATRICES

$(1$
(1) 2
(1) 2.
(1) • (4)

- 3 . 5


# "EMBEDDING" OF TREES ON BINARY MATRICES 

(1)

## "EMBEDDING" OF TREES ON BINARY MATRICES



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-Root: bottom-right node

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-Left children: go up the diagonal
-Right children: go back on each row

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## "EMBEDDING" OF TREES ON BINARY MATRICES


-Root: bottom-right node
-Left children: go up the diagonal
-Right children: go back on each row
-Left children: bounce off the diagonal

## "EmbedDing" of Trees on binary matrices



Rotate the matrix by $90^{\circ}$ to obtain the Fishburn tree!

Fishburn
Trees


$\rightarrow$ Max-Decomposition
$\leftarrow$ In-order sequence

Path label $\leftrightarrow$ Row index Node label $\leftrightarrow$ Column index

$\rightarrow$ Max-Decomposition
$\leftarrow$ In-order sequence

Path label $\leftrightarrow$ Row index Node label $\leftrightarrow$ Column index

## ApPLICATIONS: FLIP AND SUM

## FLIP AND SUM

- Duality (flip) acts as an involution on Fishburn posets.
- On Fishburn matrices, the flip corresponds to the reflection of a matrix in its antidiagonal.
- The sum of two Fishburn matrices is a Fishburn matrix.


## ApPLICATIONS: FLIP AND SUM

## FLIP AND SUM

- Duality (flip) acts as an involution on Fishburn posets.
- On Fishburn matrices, the flip corresponds to the reflection of a matrix in its antidiagonal.
- The sum of two Fishburn matrices is a Fishburn matrix.
(1) How do flip and sum act on the corresponding ascent sequences?
(2) How do flip and sum act on Fishburn trees?
(3) Is there a natural involution on the set of Fishburn trees?


## Applications: Subfamilies

Fishburn trees Modified seq. Fishburn mat. (2+2)-free posets

| Strictly decreasing | Primitive | Binary | Primitive |
| :--- | :--- | :--- | :--- |
| Comb-shaped | Self-modified | Positive diagonal | $\exists$ Chain of max length |
| $(\star)$ | $\hat{\mathcal{A}}(212,312)$ | NW-free | $N$-free (Series-parallel) |
| $(\dagger)$ | $\hat{\mathcal{A}}(231)$ | SW-free | $(\mathbf{3}+\mathbf{1})$-free (Semiorders) |
| By intersection of <br> the above two rows | $\hat{\mathcal{A}}(212,312,231)$ | $($ NW,SW)-free | $N$ - and (3+1)-free |

$$
\begin{aligned}
& (\star) \nexists u, v: \mathfrak{l}(u)<\mathfrak{l}(v) \leq \mathfrak{b}(u)<\mathfrak{b}(v) \\
& (\dagger) \nexists u, v: \mathfrak{b}(u)<\mathfrak{b}(v), \mathfrak{l}(u)>\mathfrak{l}(v)
\end{aligned}
$$

## Applications: Statistics

Fishburn trees Modified seq. Fishburn mat. (2+2)-free posets

| \# nodes | Length | Sum of entries | \# elements |
| :--- | :--- | :--- | :--- |
| Biggest node label | Maximum value | \# rows/cols | \# levels |
| $\|\{u: \mathfrak{l}(u)=j\}\|$ | \# copies of $j$ | $\sum_{i} a_{i, j}$ | $\left\|L_{j}\right\|$ |
| $\|\{u: \mathfrak{b}(u)=\mathfrak{l}(u)\}\|$ | \# weak ltr-max | Trace | $\sum_{i}\left\|L_{i} \cap D_{i+1}\right\|$ |
| $\|\{u: \mathfrak{l}(u)=1\}\|$ | \# copies of 1 | $\sum_{i} a_{i, 1}$ | $\#$ minimal elements |
| $\|\operatorname{rpath}(r)\|$ | \# weak rtl-max | $\sum_{j} a_{k, j}$ | \# maximal elements |
| $\mathfrak{l}\left(v_{n}\right)$ | $x_{n}$ | index | Minimal level of a <br> maximal element |

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Fishburn trees Modified seq. Fishburn mat. (2+2)-free posets

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## Thanks!

