OCCURRENCES OF SUBSEQUENCES IN BINARY WORDS

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 $B_{n,p}(k)$

 $c_p(w) =$ Number of subsequences of w that match p.

Example. For p = 10 and w = 010010, $c_p(w) = 4$.

010010

 $B_{n,p}(k) = \#\{w \in \{\mathbf{0}, \mathbf{1}\}^n \mid c_p(w) = k\}$

Number of binary words of length n that contain p exactly k times.

Example. If p = 1, then

$$B_{n,p}(k) = \binom{n}{k}$$

Runs

Maximal subsequence of consecutive letters that are equal.

Example. Can encode word using run lengths and first letter.

10011100

$egin{aligned} egin{aligned} egin{aligned} egin{aligned} B_{n,p}(\mathbf{0}) \ \end{array} \end{aligned}$

Idea: How much of p can we get?

Suppose p = 10010

00....0111....1

00...0

 $0 \ 0 \ \cdots \ 0 \ 1 \ 1 \ 1 \ \cdots \ 1 \ 0 \ 1 \ 1 \ \cdots \ 1$

 $(1-p_1)^{i_1} p_1 (1-p_2)^{i_2} p_2 \cdots (1-p_j)^{i_j} p_j (1-p_{j+1})^{i_{j+1}}$

Exercise. If p is of length l with r_i runs of size i, then for any $k \ge 2$, $B_{l+1,p}(k) = r_{k-1}$.

 $B_{n,p}(1)$

If p is of length l and has r runs, $B_{n,p}(1) = \binom{n-r+1}{l-r+1}.$

Idea: Insert letters into *p* without creating new occurrences.

Suppose p = 100011

 $\bigcup_{0} 1 \bigcup_{1} 0 \bigcup_{1} 0 \bigcup_{1} 1 \bigcup_{0} 1 \bigcup_{0} 1 \bigcup_{0} 0$

 $(1, 2, 3, 2)_1$

 $B_{n,p}(2)$

If p has length l and r runs with r_1 of size 1, $B_{n,p}(2) = r_1 \binom{n-r}{l-r+1}.$

Idea: Use runs of size 1 to make an extra occurrence. Insert letters without creating any more.

Suppose p = 110100

 $\begin{array}{c}
1 \ 1 \ 0 \ 1 \ 0 \ 0 \\
\downarrow \\
1 \ 1 \ 0 \ 0 \ 1 \ 0 \ 0
\end{array}$

One way to insert letters is

 $\bigcup_{0}^{0} 1 \bigcup_{1}^{0} \bigcup_{1}^{0} \bigcup_{1}^{0} \bigcup_{1}^{0} \bigcup_{1}^{0} \bigcup_{0}^{1} \bigcup_{0}^{0} \bigcup_{0}^{1} \bigcup_{0}^{0} \bigcup_$



Let p be a pattern of length l with at most 3 runs, say $p = 1^i 0^j 1^k$ where $i, k \ge 0, j \ge 1$.

Maximum Occurrences

• $M_{n,p} = \max\{c_p(w) \mid w \in \{0, 1\}^n\}.$ • $w \in \{0, 1\}^n$ is *p*-optimal if $c_p(w) = M_{n,p}.$

1. We have $M_{n,p} = \max\left\{\binom{a}{i}\binom{b}{j}\binom{c}{k} \mid a+b+c=n\right\}$.

2. If $1^a 0^b 1^c$ is *p*-optimal, then so is at least one word in $\{1^{a+1}0^b 1^c, 1^a 0^{b+1}1^c, 1^a 0^b 1^{c+1}\}.$

3. The sequence $(M_{n,p})_{n\geq 0}$ is log-concave.

Internal Zeroes

We say p has an *internal zero at* n if there exist $0 \le k_1 < k_2 < k_3$ such that

 $B_{n,p}(k_1), B_{n,p}(k_3) \neq 0$ but $B_{n,p}(k_2) = 0.$

1. p = 10 does not have any internal zeroes.

2. p = 101 does not have an internal zero at n for $n \ge 7$.

3. If p isn't alternating, it has internal zeroes at all $n \ge l+3$.

Future Directions

• Refine $B_{n,p}(k)$ by keeping track of statistics. Number of 1s, number of runs etc.

• We have used runs to study $B_{n,p}(k)$. Are there better ways to do this? ConjecturePatterns of length
at most 13 For two patterns p, q, we have $B_{n,p}(k) = B_{n,q}(k)$ for all $n, k \ge 0$ if and only if $p = q, q^r, q^c$, or q^{rc} .

Expressions for $B_{n,p}(3)$ and $B_{n,p}(4)$ in [2]. Formulas for $B_{n,p}(k)$ for $k \ge 5$? Might be easier to study for particular patterns p. Interesting case: Alternating words.

References

[1] M. Bóna, B. E. Sagan, and V. R. Vatter. Pattern frequency sequences and internal zeros. *Advances in Applied Mathematics*, 2002.
 [2] K. Menon and A. Singh. Subsequence frequency in binary words. *arXiv:2306.07870*, 2023.

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