# Generating trees and pattern-avoiding inversion sequences

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based on joint work with Toufik Mansour

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- Generating trees and kernel method
- Pattern-avoiding inversion sequences
- Some new enumerative results for inversion sequences avoiding 021 and a five-letter pattern.

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Any set C of discrete objects with a notion of a size such that for each n there are finitely many objects of size n is called a combinatorial class.

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Any set C of discrete objects with a notion of a size such that for each n there are finitely many objects of size n is called a combinatorial class.

A generating tree for C is a rooted, labeled, ordered tree whose vertices are the objects of C with the following properties:

(i) each object of C appears exactly once in the tree;

(ii) objects of size *n* appears at the level *n* in the tree;

(iii) each object's children are obtained by a set of succession rules which determines the number of children and their labels.

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Fibonacci Tree:

Root: (1), Rules: (1)  $\rightsquigarrow$  (2) (2)  $\rightsquigarrow$  (1)(2)

Catalan Tree:

Root: (1),  
Rules: 
$$(m) \rightsquigarrow (2)(3) \cdots (m+1)$$

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Fibonacci Tree:

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Root: (1),
Rules: (1) \rightsquigarrow (2)
(2) \rightsquigarrow (1)(2)
```

Let  $a_n =$  the number of vertices at the level  $n, n \ge 1$ . The root is at the level 1.

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Let  $A_i(x)$  denote the generating function corresponding to the (sub)-tree with root label (*i*). Then, we get

$$A_1(x) = x + xA_2(x)$$
  
 $A_2(x) = x + xA_1(x) + xA_2(x)$ 

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Solving these equations yields  $A_1(x) = \frac{x}{1-x-x^2}$ .

The number of vertices at each level are  $1, 1, 2, 3, 5, 8, \dots = 3$ 

Root: (1), Rules:  $(m) \rightsquigarrow (2)(3) \cdots (m+1)$ 

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Root: (1), Rules:  $(m) \rightsquigarrow (2)(3) \cdots (m+1)$  $A_m(x) = x + x[A_2(x) + A_3(x) + \cdots + A_{m+1}(x)]$  for any  $m \ge 1$ .

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$$\begin{array}{l} \text{Root: (1),}\\ \text{Rules: }(m) \rightsquigarrow (2)(3) \cdots (m+1)\\\\ A_m(x) = x + x[A_2(x) + A_3(x) + \cdots + A_{m+1}(x)] \quad \text{ for any } m \geq 1. \end{array}$$

Hence

$$A_{m+1}(x) - A_m(x) = xA_{m+2}(x).$$
  
Let  $A(x, u) = \sum_{m=1}^{\infty} A_m(x)u^{m-1}.$ 

We want to determine

$$A(x,0)=A_1(x).$$

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#### Kernel Method

Multiplying  $A_{m+1}(x) - A_m(x) = xA_{m+2}(x)$  by  $u^{m-1}$  and summing over  $m \ge 1$ , we get

$$\frac{1}{u}(A(x,u) - A_1(x)) - A(x,u) = \frac{x}{u^2}(A(x,u) - A_1(x) - uA_2(x))$$

$$(\frac{1}{u} - 1 - \frac{x}{u^2})A(x, u) = -\frac{x}{u^2}A_1(x) + \frac{1}{u}$$

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If we choose  $u = \frac{1-\sqrt{1-4x}}{2}$ , then we get  $A_1(x) = \frac{1-\sqrt{1-4x}}{2x}$ , the generating function of the Catalan numbers.

The number of vertices at each level are  $1, 1, 2, 5, 14, 42, \ldots$ 

An inversion sequence of length n is an integer sequence  $e = e_1 \cdots e_n$  such that  $0 \le e_i < i$  for each  $0 \le i \le n$ .

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An inversion sequence of length n is an integer sequence  $e = e_1 \cdots e_n$  such that  $0 \le e_i < i$  for each  $0 \le i \le n$ .

We use  $I_n$  to denote the set of inversion sequences of length n.

$$\begin{aligned} &l_2 &= \{00, 01\} \\ &l_3 &= \{000, 001, 010, 011, 002, 012\} \end{aligned}$$

There is a bijection between  $I_n$  and  $S_n$ , the set of permutations of length n.

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Patterns: words over the alphabet  $[k] := \{0, 1, \dots, k-1\}$ .

A pattern  $\tau$  is a word of length k over the alphabet [k].

There are basically thirteen patterns of length three up to order isomorphism.

 $\mathcal{P}_3 = \{000, 001, 010, 100, 011, 101, 110, 021, 012, 102, 120, 201, 210\}$ 

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An inversion sequence  $e \in I_n$  contains the pattern  $\tau$  if there is a subsequence of length k in e that is order isomorphic to  $\tau$ ; otherwise, e avoids the pattern  $\tau$ .

 $I_n(\tau)$  denotes the set of all  $\tau$ -avoiding inversion sequences of length n.

For a given set of patterns *B*, we set  $I_n(B) = \bigcap_{\tau \in B} I_n(\tau)$ .

 $e = 01102321 \in I_8$  avoids the pattern 0000 because there is no subsequence  $e_i e_j e_k e_l$  of length four in e with i < j < k < l and  $e_i = e_j = e_k = e_l$ .

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 $e = 01102321 \in I_8$  avoids the pattern 0000 because there is no subsequence  $e_i e_j e_k e_l$  of length four in e with i < j < k < l and  $e_i = e_j = e_k = e_l$ .

On the other hand, e = 01102321 contains the patterns

- 010 because it has subsequences -1 - 3 1 or - - - 232- order isomorphic to 010,
- 000 because it has subsequence -11 - 1 order isomorphic to 000.

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#### Pattern-avoiding inversion sequences

They provide a unifying interpretation that relates a vast array of combinatorial structures.

- Fibonacci numbers
- Catalan numbers
- Schröder numbers
- Euler up/down numbers
- Bell numbers

- ....

appear as enumerating sequences for pattern-restricted inversion sequences.

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# Avoiding a single pattern of length three Mansour-Shattuck(2015) and Corteel et.al. (2015)

There are basically thirteen patterns of length three up to order isomorphism

 $\mathcal{P}_3 = \{000, 001, 010, 100, 011, 101, 110, 021, 012, 102, 120, 201, 210\}$ 

►  $|I_n(012)| = F_{2n-1}$ , odd Fibonacci numbers.  $F_0 = 0, F_1 = 1, F_n = F_{n-1} + F_{n-2}$  for  $n \ge 2$ .

▶  $|I_n(021)| = r_{n-1}$ ,  $r_n$  is the  $n^{th}$  large Schröder numbers.  $r_n$  is the number of paths in the first quadrant from (0,0) to (2n,0) with steps (1,1), (1,-1) and (2,0).

►  $|I_n(000)| = E_{n+1}$  is the Euler up/down numbers.  $E_n$  is the number of permutations of length *n* such that  $\sigma_1 < \sigma_2 > \sigma_3 < \cdots$ .  $\tan x + \sec x = \sum_{n=0}^{\infty} E_n \frac{x^n}{n!}$ 

#### Avoiding a single pattern of length three

► 
$$|I_n(001)| = 2^{n-1}$$

- II<sub>n</sub>(011)| = B<sub>n</sub>, the n<sup>th</sup> Bell number.

   B<sub>n</sub> is the number of ways to partition an n−element set into non-empty subsets called blocks.
- ►  $|I_n(101)| = |I_n(110)|$  = powered Catalan numbers.
- ►  $|I_n(201)| = |I_n(210)|$  corresponds to A263777 in OEIS.
- *I<sub>n</sub>*(102)|, *I<sub>n</sub>*(010)|, *I<sub>n</sub>*(100)|, *I<sub>n</sub>*(120)| corresponds to A200753, A263779, A263780, A263778, respectively.
   The case 010 is solved by B. Testart(2022).

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Consider inversion sequences avoiding two element subsets of

 $\mathcal{P}_3 = \{000, 001, 010, 100, 011, 101, 110, 021, 012, 102, 120, 201, 210\}$ 

There are 78 pairs but Yan and Lin (2021) showed that there 48 Wilf classes and enumerated 17 of them.

$$|I_n(001, 210)| = \binom{n}{3} + n$$

- For  $p \in \{(012, 201), (012, 210)\}, |I_n(p)| = 2^{n+1} - \binom{n+1}{3} - 2n - 1$ 

- For  $p \in \{(012, 021), (110, 012)\}$ ,  $|I_n(p)| = 2^n - n$ 

## Enumeration of $I_n(B)$

algorithm+generating tree+kernel method  $\rightarrow$  enumeration

Generating function  $R(x) = \sum_{n \ge 1} |I_n(B)| x^n$ .

We have three main steps:

- determine the succession rules for the corresponding generating tree.
- determine the equations satisfied by the generating functions from the tree rules.
- solve the equations and determine the enumerating sequence for  $I_n(B)$ .

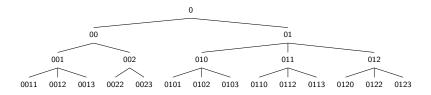
We will construct a pattern-avoidance tree  $\mathcal{T}(B)$  corresponding to the class  $\mathcal{I}(B)$  as follows:

- the root is 0 and stays at level 1.
- ▶ the  $n^{th}$  level of the tree consist exactly the elements of  $I_n(B)$ .
- ► the children of e<sub>1</sub>e<sub>2</sub>···e<sub>n-1</sub> ∈ I<sub>n-1</sub>(B) are obtained from the set {e<sub>1</sub>e<sub>2</sub>···e<sub>n-1</sub>e<sub>n</sub>|e<sub>n</sub> = 0, 1, ..., n 1} by obeying the pattern avoidance restrictions of B.
- ▶ we arrange the nodes from the left to the right so that e<sub>1</sub>e<sub>2</sub> ··· e<sub>n-1</sub>i appears on the left of e<sub>1</sub>e<sub>2</sub> ··· e<sub>n-1</sub>j if i < j.</p>

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#### Tree representation of $I_n(000, 021)$

#### First four levels of $\mathcal{T}(000, 021)$



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Let  $\mathcal{T}(B; e)$  denote the subtree in  $\mathcal{T}(B)$  which has the root e.

We define an equivalence relation on  $\mathcal{T}(B)$  as follows: let  $e, e' \in \mathcal{T}(B)$ 

 $e \sim e'$  if and only if  $\mathcal{T}(B; e) \cong \mathcal{T}(B; e')$ 

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$$e\sim e'$$
 if and only if  $\mathcal{T}(B;e)\cong \mathcal{T}(B;e')$ 

#### Lemma

Let t be the length of the longest pattern in B.

 $\mathcal{T}(B; e) \cong \mathcal{T}(B; e')$  if and only if  $\mathcal{T}^{2t}(B; e) \cong \mathcal{T}^{2t}(B; e')$ 

#### a simple example: enumeration of $I_n(000, 001, 012)$



The generating tree succession rules are given by

Root: 0, Rules:  $0 \rightsquigarrow 00, 01$  $01 \rightsquigarrow 00, 011$  $011 \rightsquigarrow 00.$ 

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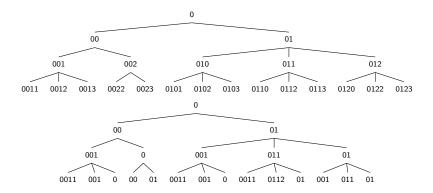
#### a simple example: enumeration of $I_n(000, 001, 012)$

Let 
$$R(x) = \sum_{n>1} |I_n(000, 001, 012)| x^n$$
.

$$\begin{aligned} R(x) &= x + xA_{00}(x) + xA_{01}(x), \\ A_{01}(x) &= x + xA_{00}(x) + xA_{011}(x), \\ A_{011}(x) &= x + xA_{00}(x), \\ A_{00}(x) &= x. \end{aligned}$$

Then  $R(x) = x^4 + 2x^3 + 2x^2 + x$ .

Based on the second tree, we try to figure out the succession rules of the generating tree for the class  $\mathcal{I}(B)$ .



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#### Succession rules for the generating tree of $I_n(000, 021)$

We define  $r_0 = 0$ ,  $b_0 = 0$ , and for  $m \ge 1$ ,

$$a_m = 0011 \cdots mm$$
  
 $b_m = 0011 \cdots (m-1)(m-1)m$   
 $c_m = 01122 \cdots mm$   
 $d_m = 01122 \cdots (m-1)(m-1)m.$ 

The generating tree  $\mathcal{T}(000, 021)$  is given by

Root: 
$$r_0$$
,  
Rules:  $r_0 \rightsquigarrow a_0 d_1$ ,  
 $a_m \rightsquigarrow b_{m+1} b_m \cdots b_0$ ,  
 $b_m \rightsquigarrow a_m b_m b_{m-1} \cdots b_0$ ,  
 $c_m \rightsquigarrow a_m d_{m+1} d_m \cdots d_1$ ,  
 $d_m \rightsquigarrow b_m c_m d_m d_{m-1} \cdots d_1$ .

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#### Equations for the generating functions

Let 
$$R(x) = \sum_{n \ge 1} |I_n(000, 021)| x^n$$
.

$$R(x) = x + xA_0(x) + xD_1(x),$$
  

$$A_m(x) = x + x \sum_{j=0}^{m+1} B_j(x),$$
  

$$B_m(x) = x + xA_m(x) + x \sum_{j=0}^{m} B_j(x),$$
  

$$C_m(x) = x + xA_m(x) + x \sum_{j=1}^{m+1} D_j(x),$$
  

$$D_m(x) = x + xB_m(x) + xC_m(x) + x \sum_{j=1}^{m} D_j(x).$$

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### Bivariate generating functions

We define

$$A(x, u) = \sum_{m \ge 0} A_m(x) u^m, B(x, u) = \sum_{m \ge 0} B_m(x) u^m$$
$$C(x, u) = \sum_{m \ge 1} C_m(x) u^{m-1}, D(x, u) = \sum_{m \ge 1} D_m(x) u^{m-1}.$$

Hence

$$\begin{aligned} R(x) &= x + xA(x,0) + xD(x,0), \\ A(x,u) &= \frac{x}{1-u} + \frac{x}{u}(B(x,u) - B(x,0)) + \frac{x}{1-u}B(x,u), \\ B(x,u) &= \frac{x}{1-u} + xA(x,u) + \frac{x}{1-u}B(x,u), \\ C(x,u) &= \frac{x}{1-u} + \frac{x}{u}(A(x,u) + D(x,u) - A(x,0) - D(x,0)) + \frac{x}{1-u}D(x,u) \\ D(x,u) &= \frac{x}{1-u} + \frac{x}{u}(B(x,u) - B(x,0)) + xC(x,u) + \frac{x}{1-u}D(x,u). \end{aligned}$$

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## Kernel Method

From the second and third equations, we get

$$\frac{(1-x)u - x^2 - u^2}{u(1-u-x)}A(x,u) = -\frac{x^2}{u(1-x)}A(x,0) + \frac{x}{(1-u-x)(1-x)}$$

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By choosing  $u = x^2 M(x)$ , where  $M(x) = \frac{1-x-\sqrt{1-2x-3x^2}}{2x^2}$  is the generating function for the Motzkin numbers, we obtain

$$A(x,0) = \frac{xM(x)}{1 - x - x^2M(x)} = xM^2(x),$$

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and then

$$A(x, u) = \frac{xM(x)(x^2M(x) - u)}{u^2 + (x - 1)u + x^2}$$

and

$$B(x, u) = \frac{xM(x)((u+x)x^2M(x) - u + x^2)}{u^2 + u(x-1) + x^2}$$

The generating function  $R(x) = \sum_{n \ge 1} |I_n(000, 021)| x^n$  is given by

$$R(x) = \frac{3x^3 + x^2 - 3x + 1}{2x^2\sqrt{(1+x)(1-3x)}} - \frac{(1-x)^2}{2x^2}.$$

Hence

$$|I_n(000,021)| = \frac{1}{2}(3a_{n-1} + a_n - 3a_{n+1} + a_{n+2})$$

for all  $n \ge 1$  where  $a_n = \sum_{k=0}^n (-1)^{n-k} {n \choose k} {2k \choose k}$ .

### Enumeration of $I_n(021, 00011)$

the vertex labeling symbols:  $a_m = 0^m, b_m = 01^m, c_m = 001^m, d_m = 0^{2}1^{2} \dots m^{2}, e_m = 0^{2}1^{2} \dots (m-1)^{2}m, f_m = 01^{2}2^{2} \dots m^{2}, and$   $g_m = 01^{2}2^{2} \dots (m-1)^{2}m, \text{ for all } m \ge 1; and a_e = e \text{ for an inversion sequence } e.$ the generating tree succession rules:

#### Then we have

$$F_{(021,00011)}(x) = \frac{(1-x-3x^2)\sqrt{1-4x}-9x^3-3x^2+4x-1}{2x^3\sqrt{(1+x)(1-3x)}} - \frac{(2-2x+3x^2)\sqrt{1-4x}+2x^3-8x^2+7x-2}{2x^3}.$$

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### Enumeration of $I_n(021, 00012)$

the vertex labeling symbols:  $a_m = 0^m, b_m = 01^m, c_m = 001^m, d_m = 0^{2}1^{2} \dots m^{2}, e_m = 0^{2}1^{2} \dots (m-1)^{2}m, f_m = 01^{2}2^{2} \dots m^{2}, and$   $g_m = 01^{2}2^{2} \dots (m-1)^{2}m, \text{ for all } m \ge 1; and a_e = e \text{ for an inversion sequence } e.$ the generating tree succession rules:

Then we have

$$F_{(021,00012)}(x) = \frac{4x^9 - 8x^8 - 8x^7 + 10x^6 - 18x^5 + 6x^4 + 8x^3 - 9x^2 + 4x - 1}{2x^2(1+x)^2(1-x)^4(1-2x)} + \frac{18x^7 - 24x^6 + 24x^5 + 8x^4 - 26x^3 + 17x^2 - 6x + 1}{2x^2(1+x)(1-x)^4(1-2x)\sqrt{(1+x)(1-3x)}}.$$

#### Enumeration of $I_n(021, \tau)$

Mansour-Y. (2022)

We determined the generating trees and generating functions for the inversion sequences avoiding 021 and another pattern of length 4 or 5.

$$|I_n(\{021,0001\})| = \frac{(4n-25)(-1)^n}{32} - \frac{n(n+1)-1}{4} + \frac{1}{32}3^{n+4} + \sum_{j=0}^{n+1} \left(\frac{(4j-39)(-1)^j}{32} + \frac{1}{4}j^2 - j + \frac{1}{2} - \frac{1}{32}3^{j+2}\right) M_{n+1-j},$$

$$|I_n(\{021,0010\})| = \binom{2n}{n},$$

$$|I_n(\{021,0011\})| = C_{n+2} + 1 - \sum_{j=0}^{n+1} C_j,$$

$$|I_n(\{021,0012\})| = 2^{n+3} - \frac{(n+1)(2n^2+7n+24)}{6} - 3,$$

$$|I_n(\{021,0100\})| = |I_n(\{021,0110\})| = \frac{n^2+n+6}{8(2n+3)(2n+5)} \binom{2n+6}{n+3},$$

$$|I_n(\{021,0101\})| = |I_n(\{021,0111\})| = \sum_{i=1}^{n+1} \frac{1}{i} \binom{n}{i-1} \binom{2n+2-i}{i-1}$$

$$|I_n(\{021,0102\})| = 2^{n+1} - \frac{(n+1)(n^2+2n+12)}{6} - 1 + \sum_{j=0}^{n+1} C_j$$

$$|I_n(\{021, 0112\})| = C_{n+1} - 2^{n+1} + 1 + \sum_{j=0}^n 2^{n-j}C_j,$$

$$|I_n(\{021,0120\})| = |I_n(\{021,0122\})| = \frac{1}{2} \binom{2n+2}{n+1} - \frac{1}{2} \sum_{j=1}^n \binom{2j}{j},$$

$$|I_n(\{021,0123\})| = 2^{n-1}(n^2 - 3n + 4) + \frac{n(n+1)}{2} - 1,$$

$$|I_n(\{021, 1000\})| = |I_n(\{021, 1100\})| = \frac{n^5 + 2n^4 + 23n^3 + 46n^2 + 120n + 48}{2(n+1)(n+2)(n+3)(n+4)} \binom{2n}{n},$$

$$|I_n(\{021, 1001\})| = |I_n(\{021, 1011\})| = |I_n(\{021, 1101\})| = \frac{1}{n+1} \sum_{j=0}^{\lfloor \frac{n}{2} \rfloor} {n+1 \choose j} {2n+2 \choose n-2j},$$

$$|I_n(\{021, 1002\})| = \frac{1}{2} \binom{2n+6}{n+3} - \frac{5}{2} \binom{2n+4}{n+2} + \frac{5}{2} \binom{2n+2}{n+1} + \frac{1}{2} \sum_{j=0}^n \binom{2j}{j}$$

$$+2^{n+1}-\frac{1}{24}(n^4+2n^3+11n^2+34n+36),$$

 $|I_n(\{021, 1020\})| = |I_n(\{021, 1022\})|$ 

$$=\binom{2n+8}{n+4}-\frac{13}{2}\binom{2n+6}{n+3}+\frac{21}{2}\binom{2n+3}{n+2}-\frac{1}{2}\sum_{j=0}^{n+1}\binom{2j}{j}-\frac{1}{2},$$

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$$\begin{split} |I_n(\{021, 1023\})| &= \sum_{j=0}^{n+1} (2^{j+1} - j - 1)C_{n+1-j} + \frac{n(3n^3 + 22n^2 + 129n + 398)}{24} + 2^{n-1}(n^2 - 3n - 52) + 24, \\ |I_n(\{021, 1102\})| &= \frac{1}{2} \binom{2n+6}{n+3} - \frac{21}{4} \binom{2n+4}{n+2} + \binom{2n+2}{n+1} + \frac{(n+1)^2}{2} - 2^n + \frac{1}{2} \sum_{j=1}^{n+3} (2^{j-2} - 3j + 8) \binom{2n+6-2j}{n+3-j} \\ |I_n(\{021, 1120\})| &= 4^n - \frac{n}{2(2n+3)} \binom{2n+4}{n+2}, \\ |I_n(\{021, 1200\})| &= |I_n(\{021, 1220\})| = \frac{n+4}{2(n+2)} \binom{2n+2}{n+1} + \sum_{j=0}^{n-1} (2j+1) \binom{2j}{j} - 4^n, \\ |I_n(\{021, 1203\})| &= \frac{n+1}{24} (n^3 + n^2 - 2 * n - 108) + 2^{n-1}(n^2 - 11n + 28) - \frac{19}{2} \\ &+ \frac{1}{2} \binom{2n+2}{n+1} - \frac{1}{2} \sum_{j=2}^{n+1} (j - 1) \binom{2n+2-2j}{n+1-j}, \\ |I_n(\{021, 1220\})| &= \frac{n+4}{n+2} \binom{2n+1}{n} - 4^n + \sum_{j=0}^{n} (2j+1) \binom{2j}{j}, \\ |I_n(\{021, 1230\})| &= \frac{1}{3} (2 \cdot 4^n + 1) + \sum_{j=1}^{n} (\frac{j}{2} - 2^{j-1}) \binom{2n+2-2j}{n+1-j}. \end{split}$$

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# Thank you!

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#### succession rules

Consider the children of a node labeled by  $a_m = 0011 \cdots mm$ :

 $a_m j = 0011 \cdots mm j$  where  $j = m + 1, m + 2, \dots, 2m + 2$ otherwise,  $a_m j$  does not avoid B.

• 
$$a_m(m+1) = 0011 \cdots mm(m+1) = b_{m+1};$$

- For other j values, a<sub>m</sub>(m + j) = 0011 · · · mm(m + j); note that T(B; a<sub>m</sub>(m + j)) ≅ T(B; b<sub>m+2-j</sub>) by removing the letters m + 2 − j, m + 3 − j, . . . , m and decreasing each letter greater than m by 2j − 1.
- ► therefore the children of the node with label a<sub>m</sub> are exactly the nodes labelled by b<sub>m+1</sub>, b<sub>m</sub>,..., b<sub>0</sub>, that is, a<sub>m</sub> ~→ b<sub>m+1</sub>b<sub>m</sub> ··· b<sub>0</sub>.